## Oddtown/Eventown

Eventown: $\mathcal{F} \subseteq[n]$ is an Eventown-family of sets if

- $|F| \equiv 0(\bmod 2)$ for all $F \in \mathcal{F}$ and
- $\left|F_{1} \cap F_{2}\right| \equiv 0(\bmod 2)$ for every $F_{1}, F_{2} \in \mathcal{F}$

How large can $|\mathcal{F}|$ be? As large as $2\lfloor n / 2\rfloor$

Construction. For even $n$ :
$\mathcal{F}=\left\{F \subseteq[n]:|F \cap\{2 i-1,2 i\}|\right.$ is even for all $\left.i \in\left[\frac{n}{2}\right]\right\}$

Oddtown: $\mathcal{F} \subseteq[n]$ is an Oddtown-family of sets if

- $|F| \equiv 1(\bmod 2)$ for all $F \in \mathcal{F}$ and
- $\left|F_{1} \cap F_{2}\right| \equiv 0(\bmod 2)$ for every $F_{1} \neq F_{2} \in \mathcal{F}$

How large can $|\mathcal{F}|$ be?

Oddtown Theorem The maximum size of an Oddtownfamily over [ $n$ ] is $n$.

Proof. Let $\mathcal{F}=\left\{F_{1}, \ldots, F_{m}\right\} \subseteq 2^{[n]}$ be an Oddtownfamily.

Let $\mathbf{v}_{\mathbf{i}} \in\{0,1\}^{n}$ be the characteristic vector of $F_{i}$ : $j^{\text {th }}$ coordinate is 1 if $j \in F_{i}$, otherwise 0 .

Crucial property: $\mathbf{v}_{\mathbf{i}}{ }^{T} \mathbf{v}_{\mathbf{j}}=\left|F_{i} \cap F_{j}\right|$
Claim $\mathbf{v}_{\mathbf{1}}, \ldots, \mathrm{v}_{\mathbf{n}}$ is linearly independent over $\mathbb{F}_{2}$.
Let $\lambda_{1} \mathbf{v}_{\mathbf{1}}+\cdots+\lambda_{m} \mathbf{v}_{\mathbf{m}}=0$
Then for every $i$

$$
\begin{aligned}
0 & =\left(\lambda_{1} \mathbf{v}_{1}+\cdots+\lambda_{m} \mathbf{v}_{\mathbf{m}}\right)^{T} \mathbf{v}_{\mathbf{i}} \\
& =\lambda_{1} \mathbf{v}_{\mathbf{1}}^{T} \mathbf{v}_{\mathbf{i}}+\cdots+\lambda_{i} \mathbf{v}_{\mathbf{i}}{ }^{T} \mathbf{v}_{\mathbf{i}}+\cdots \lambda_{m} \mathbf{v}_{\mathbf{m}}{ }^{T} \mathbf{v}_{\mathbf{i}} \\
& =\lambda_{i}
\end{aligned}
$$

Since $\mathbf{v}_{\mathbf{1}}, \ldots \mathbf{v}_{\mathbf{m}}$ are linearly independent vectors in an $n$-dimensional space, $m \leq n$.

## The Linear Algebra bound

The key to the Oddtown proof is the following simple observation:

Linear Algebra bound If $\mathbf{v}_{\mathbf{1}}, \ldots, \mathbf{v}_{\mathbf{m}}$ are a set of linearly independent vectors belonging to the span of the vectors $\mathbf{u}_{1}, \ldots \mathbf{u}_{\mathbf{k}}$, then $m \leq k$.

The Gram matrix $M=$ ( $m_{i j}$ ) of a set of vectors $\mathbf{v}_{\mathbf{1}}, \ldots, \mathbf{v}_{\mathbf{m}}$ is defined by $m_{i j}=\mathbf{v}_{\mathbf{i}}{ }^{T} \mathbf{v}_{\mathbf{j}}$.
Proposition Vectors $\mathbf{v}_{\mathbf{1}}, \ldots, \mathrm{v}_{\mathrm{m}} \in \mathbb{F}^{n}$ are linearly independent iff their Gram matrix over $\mathbb{F}$ is nonsingular.

Proof of Oddtown Thm. The Gram matrix of the characteristic vectors of an Oddtown family over $\mathbb{F}_{2}$ is the identity matrix. Then Apply Linear Algebra bound. $\square$

## Explicit Ramsey graphs I

A clique or an independent set of a graph $G$ is called a homogenous set.

A graph is $k$-Ramsey if it does not contain a homogenous set of order $k$. (Remark: instead of RED/BLUEcoloring we formulate in terms of edge/non-edge.)

In a week we will know that the largest $k$-Ramsey graph has at least $\sqrt{2}^{k}$ vertices. ( $R(k, k) \geq \sqrt{2}^{k}$.) We will prove its existence by the probabilistic method.

BUT: can you give one such beast in my hand?
Sure: go over all the $2\left(\begin{array}{c}\binom{n}{2} \\ \text { graphs on } n=\sqrt{2}^{k} \text { verti- }\end{array}\right.$ ces and check whether their clique number and independence number are below $k$ (Never mind that these are NP-hard problems).
Eventually you'll find a $k$-Ramsey graph.

## Explicit Ramsey graphs II

Why are you not happy with this "construction"?
It takes too much time.
What is then a constructive $k$-Ramsey graph?
Its adjacency matrix should be constructible in time polynomial in its number of vertices $n$.
Or even stronger: adjacency of any two vertices should be decidable in time polynomial in $\log n$ (what it takes to write down the label of the two vertices).

Turán construction: $T_{(k-1)^{2}, k-1}$ has no clique and no independent set of order $k$. Has $(k-1)^{2}$ vertices.

Anything more than quadratic?

## Construction of Nagy

$V(G)=\binom{[k]}{3}$,
$E(G)=\{A B:|A \cap B|=1\}$
For $k=3 \leadsto$ Petersen graph
Theorem $G$ has no homogenous set of order $k+1$. Proof. An independent set of $G$ is an Oddtown family $\Rightarrow \quad \alpha(G) \leq k$.
For the vertices $C_{1}, \ldots, C_{m}$ of a clique, we have that

- $\left|C_{i}\right|=3$ for all $i$ and
- $\left|C_{i} \cap C_{j}\right|=1$ for every $i \neq j$

Hence the Gram matrix of the characteristic vectors is $J_{m}+2 I_{m}$ ( $J_{m}$ is the all 1 matrix, $I_{m}$ is the identity) $M$ is nonsingular over $\mathbb{R} \Rightarrow \mathbf{v}_{\mathbf{1}}, \ldots, \mathbf{v}_{\mathbf{m}} \in R^{n}$ are linearly independent $\Rightarrow m \leq n$.
So $\omega(G) \leq n$. $\square$
$G$ is a $k$-Ramsey graph on $\Theta\left(k^{3}\right)$ vertices.
HW: Linear Algebra-free proof (like original)

Non-uniform Fischer Inequality

Non-uniform Fischer Inequality Let $\lambda \in \mathbb{N}$. If $\mathcal{F}=$ $\left\{F_{1}, \ldots, F_{m}\right\} \subseteq 2^{[n]}$ is a family of sets with $\mid C_{i} \cap$ $C_{j} \mid=\lambda$ for all $i \neq j$, then

$$
|\mathcal{F}| \leq n .
$$

Proof: Case $1 \exists i$ such that $\left|F_{i}\right|=\lambda$.
Consider $\left\{F_{j} \backslash F_{i}: j \neq i\right\}$ and use induction on $n$.

Case 2: $\forall i$, we have $\left|F_{i}\right|>\lambda$.
The Gram matrix of $\mathcal{F}$ is $\lambda J_{m}+D$, where $D$ is the diagonal matrix with diagonal entry $d_{i}=\left|F_{i}\right|-\lambda>0$, $i=1, \ldots, m$.
$\mathbf{x}^{T}(\lambda J+D) \mathbf{x}=\lambda\left(\sum_{i=1}^{m} \mathrm{x}_{i}^{2}\right)+\sum_{i=1}^{m} d_{i} \mathrm{x}_{i}^{2}>0$ for every $\mathrm{x} \neq 0$, so the Gram matrix is positive definite.
$\Rightarrow$ the Gram matrix has full rank ( $m$ ) over $\mathbb{R}$.
$\Rightarrow$ the charactersitic vectors are linearly independent $\Rightarrow m \leq n$.

## Construction of Frankl-Wilson


$E(G)=\{A B:|A \cap B| \equiv-1(\bmod p)\}$
Remark $p=2 \leadsto$ Nagy graph

Theorem (Frankl-Wilson) The graph $G$ above gives a constructive $k$-Ramsey with a vertex set size of order

$$
k^{O\left(\frac{\ln k}{\ln \ln k}\right)} .
$$

## Proof: Next semester in DM III.

## Chromatic number of the unit-distance graph

$G_{n}$ is the $n$-dimensional unit distance graph.
$V\left(G_{n}\right)=\mathbb{R}^{n}$
$E\left(G_{n}\right)=\{\mathrm{xy}:\|\mathrm{x}-\mathrm{y}\|=1\}$
\$1000 dollar question: What is the chromatic number of the plane?
We know $4 \leq \chi\left(G_{2}\right) \leq 7$. (HW)
Hadwiger-Nelson problem How fast does $\chi\left(G_{n}\right)$ grow?
Claim $\chi\left(G_{n}\right) \leq n^{n / 2}$. (HW)

$$
\chi\left(G_{n}\right) \geq n+1 \text {. (simplex with unit sidelength) }
$$

Larman-Rogers (1972) $\chi\left(G_{n}\right) \leq$ const $^{n}(\mathrm{HW})$

$$
\chi\left(G_{n}\right)=\Omega\left(n^{2}\right)
$$

Remark. Clearly, unit-distance plays no special role here. $G_{n} \cong G_{n}^{\delta}$ where $G_{n}^{\delta}$ is the " $\delta$-distance graph":
$V\left(G_{n}^{\delta}\right)=\mathbb{R}^{n}$
$E\left(G_{n}^{\delta}\right)=\{\mathbf{x y}:\|\mathrm{x}-\mathrm{y}\|=\delta\}$

## The growth of $\chi\left(G_{n}\right)$ is exponential

Theorem (Frankl-Wilson, 1981) $\chi\left(G_{n}\right) \geq \Omega\left(1.1^{n}\right)$.
Proof. Goal: For some distance $\delta>0$ we find a subgraph $H_{n} \subseteq G_{n}^{\delta}$ with $\alpha(H) \leq \frac{|V(H)|}{1.1^{n}}$.
Key: If $\mathbf{v}_{A}$ and $\mathbf{v}_{B} \in \mathbb{R}^{n}$ are the characteristic vectors of sets $A$ and $B \in 2^{[n]}$, then distance of $\mathbf{v}_{A}$ and $\mathbf{v}_{B}$ is equal to $\sqrt{|A \triangle B|}$.
If $A, B \in \mathcal{F} \subseteq\binom{[d]}{k}$ are members of a uniform family $\mathcal{F}$, then the distance of $\mathbf{v}_{A}$ and $\mathbf{v}_{B}$ depends on the intersection size: $\left\|\mathbf{v}_{A}-\mathbf{v}_{B}\right\|=\sqrt{2(k-|A \cap B|)}$.
an independent set in $G_{n}^{\delta}$ avoids distance $\delta \leadsto$ a uniform family avoiding a certain intersection size.

We give a family $\mathcal{F}$ where any subfamily $\mathcal{F}^{\prime} \subseteq \mathcal{F}$, whose members avoid a certain intersection size, is small compared to $|\mathcal{F}|$ :
Let $\mathcal{F}:=\binom{[4 p-1]}{2 p-1}$, where $p$ is a prime.
Then pairwise intersection size $p-1$ is hard to avoid.

## Generalized Oddtown

Replacing $\mathbb{F}_{2}$ in the Oddtown proof with $\mathbb{F}_{p}$ in the proof immediately gives the
modp-town Theorem. Let $p$ be a prime number and $\mathcal{F} \subseteq 2^{[n]}$ be a family such that

- $|F| \not \equiv \equiv(\bmod p)$ for all $F \in \mathcal{F}$ and
- $\left|F_{1} \cap F_{2}\right| \equiv 0(\bmod p)$ for every $F_{1} \neq F_{2} \in \mathcal{F}$

Then $|\mathcal{F}| \leq n$.

## Even More Generalized Oddtown

Theorem ("Nonuniform modular RW-Theorem", FranklWilson, 1981; Deza-Frankl-Singhi, 1983)
Let $p$ be a prime, and $L$ be a set of $s$ integers.
Let $B_{1}, \ldots, B_{m} \in 2^{[n]}$ be a family such that

- $\left|B_{i}\right| \notin L(\bmod p)$
- $\left|B_{i} \cap B_{j}\right| \in L(\bmod p)$ for every $i \neq j$.

Then

$$
m \leq \sum_{i=0}^{s}\binom{n}{i}
$$

Remark RW stands for Ray-Chaudhuri and Wilson.
Remark Oddtown Theorem: $p=2, L=\{0\}$.
The statement only gives $m \leq n+1$, but the proof will give $m \leq n$ (because $L=\{0\}$ )

## Generalizing linear independence of vectors

Let $\mathbb{F}$ be a field and $\Omega$ an arbitrary set. Then the set $\mathbb{F}^{\Omega}=\{f: \Omega \rightarrow \mathbb{F}\}$ of functions is a vector space over $\mathbb{F}$.

Lemma Let $\Omega \subseteq \mathbb{F}^{n}$. If $f_{1}, \ldots, f_{m} \in \mathbb{F}^{\Omega}$ and there exist $\mathbf{v}_{\mathbf{1}}, \ldots, \mathbf{v}_{\mathbf{m}} \in \Omega$ such that

- $f_{i}\left(\mathbf{v}_{\mathbf{i}}\right) \neq 0$, and
- $f_{i}\left(\mathbf{v}_{\mathbf{j}}\right)=0$ for all $j<i$,
then $f_{1}, \ldots, f_{m}$ are linearly independent in $\mathbb{F}^{\Omega}$.
Proof. Suppose $\lambda_{1} f_{1}+\cdots+\lambda_{m} f_{m}=0$, and let $j$ be the smallest index with $\lambda_{j} \neq 0$. Substituting $\mathbf{v}_{\mathbf{j}}$ into this function equation we have

$$
\begin{aligned}
& \underbrace{\lambda_{1} f_{1}\left(\mathbf{v}_{\mathbf{j}}\right)+\cdots+\lambda_{j-1} f_{j-1}\left(\mathbf{v}_{\mathbf{j}}\right)}_{=0, \text { since } \lambda_{i}=0, i<j}+\underbrace{\lambda_{j} f_{j}\left(\mathbf{v}_{\mathbf{j}}\right)}_{\neq 0} \\
& +\underbrace{\lambda_{j+1} f_{j+1}\left(\mathbf{v}_{\mathbf{j}}\right)+\cdots+\lambda_{m} f_{m}\left(\mathbf{v}_{\mathbf{j}}\right)}_{=0, \text { since } f_{i}\left(\mathbf{v}_{\mathbf{j}}\right)=0, j<i}=0,
\end{aligned}
$$

a contradiction.

## Proof of Even More Generalized Oddtown

For each set $B_{i}$, let $\mathbf{v}_{\mathbf{i}} \in \mathbb{F}_{p}$ be its characteristic vector. For $\mathbf{x}=\left(x_{1}, \ldots, x_{n}\right)$ let

$$
f_{i}(\mathbf{x})=\prod_{l \in L}\left(\mathbf{x}^{T} \mathbf{v}_{\mathbf{i}}-l\right)
$$

Clearly,

$$
f_{i}\left(\mathbf{v}_{\mathbf{j}}\right) \begin{cases}\neq 0 & \text { if } i=j \\ =0 & \text { if } i \neq j\end{cases}
$$

So the functions $f_{1}, \ldots, f_{m}$ are linearly independent in the subspace they generate in $\mathbb{F}_{p}\left[x_{1}, \ldots, x_{m}\right]$. What is the dimension?

Each $f_{i}$ is the product of $s$ linear functions in $n$ variables. Expanding the parenthesis: $f_{i}$ is the linear combination of terms of the form $x_{1}^{s_{1}} \cdots \cdots x_{n}^{s_{n}}$ with $s_{1}+\cdots+s_{n}=s$ ?
How many terms like that are there?
Much more than we can afford ...

## Multilinearization

We need another trick to reduce the dimension. We use that our vectors (witnessing the linear independence in the Lemma) have only 0 or 1 coordinates.

From $f_{i}$ define $\tilde{f}_{i}$ by expanding the product and replacing each power $x_{i}^{k}$ by a term $x_{i}$ for every $k \geq 1$ and $i, 1 \leq i \leq m$.

Since $0^{k}=0$ and $1^{k}=1$ for every $k \geq 1$ we have that $f_{i}\left(\mathbf{v}_{\mathbf{j}}\right)=\tilde{f}_{i}\left(\mathbf{v}_{\mathbf{j}}\right)$ for every $i, j$.

The properties of the functions and vectors remains valid, so the (now) multilinear polynomials $\tilde{f_{1}}, \ldots, \tilde{f_{m}}$ of total degree $s$ are also linearly independent.

They live in a space spanned by the basic monomials $\prod_{j=1}^{k} x_{i_{j}}$ of degree at most $s$. Their number is at most

$$
\binom{n}{s}+\binom{n}{s-1}+\cdots+\binom{n}{1}+\binom{n}{0} .
$$

## Avoiding a certain intersection size

Theorem Let $p$ be a prime number. If $\mathcal{F}^{\prime} \subseteq\binom{[4 p-1]}{2 p-1}$ such that for all $A, B \in \mathcal{F}^{\prime}$ we have $|A \cap B| \neq p-1$, then

$$
|\mathcal{F}| \leq 2 \cdot\binom{4 p-1}{p-1}<1.76^{n}
$$

Proof. Consequence of Generalized Oddtown with $L=$ $\{0,1,2, \ldots p-2\}$.

Let $n=4 p-1, k=2 p-1$.
Let $\delta=\sqrt{2(2 p-1-(p-1))}=\sqrt{2 p}$ and define $H \subseteq G_{d}^{\delta}$ by $V(H)=\left\{\mathbf{v}_{A}: A \in\binom{[n]}{k}\right\}$.
Then distance $\delta$ is hard to avoid in $V(H)$ :

$$
\alpha(H) \leq 1.76^{n}<\frac{\binom{4 p-1}{2 p-1}}{1.1^{n}} .
$$

Remark Optimizing parameters gives $\chi\left(G_{n}\right) \geq \Omega\left(1.2^{n}\right)$.

## Proof of Eventown bound

Let $\left\{F_{1}, \ldots, F_{m}\right\}$ be an Eventown family.
Let $V_{\mathcal{F}}=\left\langle\mathbf{v}_{1}, \ldots, \mathbf{v}_{m}\right\rangle \leq \mathbb{F}_{2}^{n}$ be the linear space spanned by the characteristic vectors.
Let $\mathbf{u}_{1}, \ldots, \mathbf{u}_{k} \in V_{\mathcal{F}}$ be a basis of $V_{\mathcal{F}}$.
Let $U: \mathbb{F}_{2}^{n} \rightarrow \mathbb{F}_{2}^{k}$ be the linear function defined by $U(\mathrm{x})=\left(\mathrm{x} \cdot \mathrm{u}_{1}, \ldots, \mathrm{x} \cdot \mathrm{u}_{k}\right)$.
$\mathbf{u}_{1}, \ldots, \mathbf{u}_{k}$ linearly independent, so $\operatorname{dim} \operatorname{im}(U)=k$.
Claim $V_{\mathcal{F}} \subseteq \operatorname{ker}(U)$
Proof. Any $\mathrm{x} \in V_{\mathcal{F}}$ is a linear combination of $\mathrm{v}_{i} \mathbf{s}$, so is any $\mathbf{u}_{i}$. So

$$
\mathbf{x} \cdot \mathbf{u}_{i}=\sum_{j, k} \alpha_{j} \beta_{k} \mathbf{v}_{j} \cdot \mathbf{v}_{k}=0
$$

since by the Eventown rules $\Rightarrow \mathbf{v}_{j} \cdot \mathbf{v}_{l}=0$ for every $1 \leq j, l \leq m$
$k=\operatorname{dim}\left(V_{\mathcal{F}}\right) \leq \operatorname{dim} \operatorname{ker}(U)=n-k$
$\operatorname{dim}\left(V_{\mathcal{F}}\right) \leq\left\lfloor\frac{n}{2}\right\rfloor \Rightarrow m \leq\left|V_{\mathcal{F}}\right| \leq 2^{\lfloor n / 2\rfloor}$.

