Erdős-Simonovits-Stone Theorem\_\_\_\_\_

**Theorem.** (Erdős-Stone, 1946) For arbitrary fixed integers  $r \ge 2$  and  $t \ge 1$ 

$$ex(n,T_{rt,r}) = \left(1 - \frac{1}{r-1}\right)\binom{n}{2} + o(n^2).$$

**Corollary.** (Erdős-Simonovits, 1966) For any graph H,

$$ex(n,H) = \left(1 - \frac{1}{\chi(H) - 1}\right) \binom{n}{2} + o(n^2).$$

**Corollaries of the Corollary.** 

$$ex(n, \text{octahedron}) = \frac{n^2}{4} + o(n^2)$$
$$ex(n, \text{dodecahedron}) = \frac{n^2}{4} + o(n^2)$$
$$ex(n, \text{icosahedron}) = \frac{n^2}{3} + o(n^2)$$
$$ex(n, \text{cube}) = o(n^2)$$

Proof of the Erdős-Simonovits Corollary\_

**Theorem.** (Erdős-Stone, 1946) For arbitrary fixed integers  $r \ge 2$  and  $t \ge 1$ 

$$ex(n, T_{rt,r}) = \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + o(n^2).$$

**Corollary.** (Erdős-Simonovits, 1966) For any graph H,

$$ex(n,H) = \left(1 - \frac{1}{\chi(H) - 1}\right) \binom{n}{2} + o(n^2).$$

Proof of the Corollary. Let  $r = \chi(H)$ .

- $\chi(T_{n,r-1}) < \chi(H)$ , so  $e(T_{n,r-1}) \le ex(n,H)$ .
- $T_{r\alpha,r} \supseteq H$ , so  $ex(n, T_{r\alpha,r}) \ge ex(n, H)$ , where  $\alpha$  is a constant depending on H; say  $\alpha = \alpha(H)$ .

Proof of the Erdős-Stone Thm\_

**Erdős-Stone Theorem.** (Understanding precisely what it actually says) For any  $\epsilon > 0$  and integers  $r \ge 2$ ,  $t \ge 1$  there exists an integer  $M = M(r, t, \epsilon)$ , such that any graph G on  $n \ge M$  vertices with more than  $\left(1 - \frac{1}{r-1} + \epsilon\right) {n \choose 2}$  edges contains  $T_{rt,r}$ .

We derive this through the following statement.

Seemingly Weaker Theorem. For any  $\epsilon > 0$  and integers  $r \ge 2, t \ge 1$  there exists an integer  $N = N(r, t, \epsilon)$ , such that any graph G on  $n \ge N$  vertices and with  $\delta(G) \ge \left(1 - \frac{1}{r-1} + \epsilon\right) n$  contains  $T_{rt,r}$ .

Note that w.l.o.g.  $\epsilon < \frac{1}{r-1}$ .

Derivation of the Erdős-Stone Theorem from the Seemingly Weaker Theorem.

Let *G* be a graph on  $n \ge M(r, t, \epsilon)^*$  vertices with more than  $\left(1 - \frac{1}{r-1} + \epsilon\right) {n \choose 2}$  edges. Recursively delete vertices which are adjacent to less than  $\left(1 - \frac{1}{r-1} + \frac{\epsilon}{2}\right)$ -fraction of the remaining vertices.

What is the number n' of vertices we are left with?

We deleted at most 
$$\sum_{j=n'+1}^{n} j\left(1 - \frac{1}{r-1} + \frac{\epsilon}{2}\right)$$
 ed-

ges. So

$$e(G) \leq \left(\binom{n+1}{2} - \binom{n'+1}{2}\right) \left(1 - \frac{1}{r-1} + \frac{\epsilon}{2}\right) + \binom{n'}{2}.$$

This implies

$$\frac{\epsilon}{2} \binom{n}{2} - n \le \left(\frac{1}{r-1} - \frac{\epsilon}{2}\right) \binom{n'}{2} - n'.$$

We choose  $M(r, t, \epsilon)$  such that  $n \ge M(r, t, \epsilon)$  implies  $n' \ge N(r, t, \epsilon/2)$ .

\*At this point we don't know  $M(r, t, \epsilon)$  yet!!! We'll define it in the proof through  $N(r, t, \epsilon/2)$ . (which is known!)

## Proof of the Seemingly Weaker Theorem.

Induction on r.

For r = 2 the claim is true provided  $\frac{\binom{\epsilon n}{t}n}{\binom{n}{t}} > t - 1$ , which is certainly true from some threshold  $N(2, t, \epsilon)$ .

Let  $r \ge 2$  and G be a graph on  $n \ge N(r+1, t, \epsilon)^*$ vertices with  $\delta(G) \ge \left(1 - \frac{1}{r} + \epsilon\right) n$ . We would like to find a  $T_{(r+1)t,r+1}$  in G.

Let  $s = \left\lceil \frac{t}{\epsilon} \right\rceil$ . By the induction hypothesis<sup>†</sup> there is a  $T_{rs,r}$  in G with vertex-set  $A_1 \cup \ldots \cup A_r$ , where  $|A_1| = \ldots = |A_r| = s$ .

$$U = V(G) \setminus (A_1 \cup \ldots \cup A_r).$$

 $W = \{w \in U : |N(w) \cap A_i| \ge t, i = 1, ..., r\}$ is the set of vertices eligible to extend some part of  $A_1, ..., A_r$  into a  $T_{(r+1)t,r+1}$ .

\*Again, we don't know  $N(r + 1, t, \epsilon)$  yet. <sup>†</sup>Here we assume  $N(r + 1, t, \epsilon) \ge N(r, s, \epsilon)$ . Double-count the number of edges missing between U and  $A_1 \cup \ldots \cup A_r$ . They are

- at least (|U| |W|)(s t) ( $\approx (s t)n$  if W is small)
- at most  $rs\left(rac{1}{r}-\epsilon
  ight)n$  ( $pprox (s-rt)n, rac{1}{r}$  if W is small)

From this we have

$$|W| \ge \frac{(r-1)\epsilon}{1-\epsilon}n - rs$$

Thus if n is large enough<sup>\*</sup> then

$$|W| > {\binom{s}{t}}^r (t-1).$$

So we can select t vertices from W, which are adjacent to the same t vertices in each  $A_i$ .

\*If 
$$N(r+1,t,\epsilon) > \left(\binom{s}{t}^r(t-1)+rs\right)\frac{1-\epsilon}{(r-1)\epsilon}$$